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CS514 Assignment 3

**Video Monitor: Video Distribution and Control Framework**

**Design Overview**

The Video Monitor system based on Liveobjects allows for multiple video sources to be dynamically viewed as best fits the user application.  The system consists of three primary components, a Viewer, a Server, and a Source object.  Multiple servers are interconnected with any number of Sources providing a video stream.  The servers maintain a hierarchy of services, an index to the attached sources and servers currently active.  A viewer component can then connect to a server and receive this index from which sources of interest can be selected.  The registered viewer service on the server maintains a list of sources for each viewer, the stream from that selection of sources is then pushed to that viewer.  Additionally, some sources are capable of connecting to streams with additional functionality such as resolution settings, pan-zoom-tilt control, etc.  These additional features are classified as subservices and can be utilized by the server or viewer.

[Image: Visio High level connectivity diagram]

The component roles can be classified as follows:

**VideoSource**

* Interface to an array of video sources (camera, internet based webcam, any dynamic image source, etc)
* Push video stream to connected server

|  |  |
| --- | --- |
| source_video.PNG  Figure 1 Streaming Video Source | source_web.PNG  Figure 2 Streaming Webcam Source |

**VideoServer:**

* Service Indexer
* Source Stream Distributor
* Additional features such as frame rate throttling

**VideoViewer**

* Interface which allows the user to select desired services
* Control which allows user to modify settings on the Server or Source services

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| --- | --- |
| viewer_video.PNG  Figure 3 Streaming Video Viewer | viewer_web.PNG  Figure 4 Streaming Webcam Viewer |

**VideoSource Overview**

The VideoSource takes an arbitrary video stream such as a webcam or any image source off the internet samples the source and formats the sample as a serialized bitmap.  The serialized bitmap is then pushed to the server.  Other formats were also envisioned such as an MPEG encoder / decoder.  The IP TV suggestion from Professor Birman was quite interesting although starting with a straightforward example such as a series of serialized bitmaps would be a good place to start.  Protocol contraints could more easily be classified and corrected resulting in a more easily implemented system with a base functionality set.

The VideoSource also provides an interface to special functionality provided by the actual media source.  For example, cameras such as the Logitech QuickCam Orbit MP provides a motorized pan-zoom-tilt control.  Drivers for this functionality are provided through logitech.  Additional functionality such as face tracking are also accessible.  The VideoSource classifies itself indicating these types of functionality are available.  Alternate camera systems such as Axis Web Cams provide similar functionality though the control interface on the camera is different.  The VideoSource abstracts these differences away.

The current set of functionality includes support for serialized bitmaps from either a usb based webcam or a url accessible source.  Many open webcams were easily located and serve as useful example and test input for the system.  The user interface is extremely simple providing a preview of what is sent from the source as well as essentially an on / off control to start or stop streaming.  The VideoSource connects to the source channel on startup, the details of this process are covered in the architecture section below.

**VideoServer Overview**

The VideoServer serves as a service indexer and stream distributor.  Viewer, Source and other Servers are registered allowing for information routing via other servers or directly to endpoints such as the viewer or source.  Each source is responsible for pushing frames to the source channel, these frames are captured by the server, addressed and retransmitted on the viewer channel to the viewers registered for a particular stream.

The VideoServer handles abstract frame and command objects, passing the information as necessary between channels.  Modifications to command types, video streaming approach etc do not impact how the server operates.   
  
**VideoViewer Overview**

The VideoViewer provides the user a list of available services determined through the connected VideoServer.  The user is able to view several of these streams at once currently the number is four although this is easily changed.  Viewing a stream is as simple as dragging the indexed Source to a viewing window.  The Viewer indicates to the Server that it is now interested in a particular source and frames pushed from that source will now be additionally addressed to the new viewer.

The VideoViewer also provides a custom user interface for a variety of service types.  An interface for pan-zoom-tilt control appears for sources providing this functionality.  Frame throttling can either be auto-calculated based on packet loss rates or manually controlled.  Currently the VideoViewer is functional minus the ability to issue commands. 

**Network Architecture**

**Features**

* Self-healing design to account for failures
* Robust accounting for currently available services over the channel
* Low-overhead in message passing due to relaxed near-real-time requirement of network modeling

To effectively facilitate the transport of messages and maintain the network stability, a self-healing network model is needed.  This model describes the interaction between “nodes” or LiveObjects objects interacting with the network.  To make use of the LiveObjects componentized model as well as to provide obfuscated function to objects that wish to interact with the network, the main networking component is encapsulated into a StreamProcessor class which describes the entire interaction between Nodes.  This Stream Processor class provides exported functionality to an application via an interface endpoint.

**Design**

    The self-healing network designed for this implementation takes the form of a distributed, ring-monitored multi-cast channel.  This design is used to both provide a comprehensive, near-live list of available objects connecting to the channel, maintained with an ordered check system for node failures.  The current network modeled locally in each node and ordered by node numbers.  Any nodes joining the network will be given this network model object and then broadcast their local services to be added to everyone’s model individually and resorted by node id.  This list has 2 key values: root and tail.  Root represents the node with the lowest numerical value (added longest ago), and tail, the node with the highest numerical value.  The network is designed to provide the following aggregated membership functions:

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| --- | --- |
| **Join** | When a new node joins the network, it sends out a broadcasted message reporting such and requesting a copy of the network model.  Over all the nodes, as current network model is synchronized everywhere, the tail node (highest numerical number) responds with a copy of the network model.   This new node is provisionally added to the tail node’s network model, thus blocking it from responding to any new requests.   The new node receives the model and sets his ID to the tail node’s id + 1.    The new node now sets an internal reference to the tail node as it’s “parent” node, and then transmits an ***expose*** message to the network with its service information so it can be replicated everywhere.    If a new node receives no response in request to asking for the network model, this implies one of two things.  Either the node is the root to the network; in which case it will max out its number of times it sends the request message, or one has just added and will be available soon; and thus a subsequent request should be responded to. |
| **Leave** | If a node is gracious enough to report leaving the network, it simply reports itself dead and will no-longer be considered for network duties. |
| **Expose** | To expose a service (node) to the network allows the network to know about it.  This can also be used to update the available services on a node or to replace a temporary listing in its parent.  A node broadcasts a service object that represents its local services to the channel and is integrated into every node’s network model.  Once this is completed by a new node, it is fully initialized and ready to perform network functions.  Upon any new (non-update) expose messages to the network, the root node assigns its parent reference to the new node, thus completing the ring structure of the network. |
| **Test-Live** | This is a two-way request between a node and its parent.  On a pre-set interval, a node sends this message to its parent, asking it to respond to make sure it still exists on the network.  If the message times out, it is resent a specified number of times, upon exhausting which, the parent is assumed dead.  The node then transmits a ***report-dead*** message to the network indicating that this node has gone offline and should be ignored. |
| **Report-Dead** | If a node is reported dead, a message is broadcasted to the channel, and each node removes it from its network model.    If a node is receives a message indicating itself has been reported dead, it introduces itself to the network the same way a new node would again, only changing its id if it believed itself to be the tail element before.  When a node is reported dead, each node then re-checks its parental associations.  If a node finds it has become the root node, it reorganizes accordingly and begins taking responsibility for the root’s actions. |
| **Checkup** | This consensus functionality begins with the root broadcasting a copy of it network model to the channel.  Upon receiving a checkup message, each node responds with any discrepancies it finds between the root’s network and its own, only sending services that exist locally but not in the root.  Any services existing in the root but not locally are added locally.  The root, upon receiving discrepancies, adds them to its own list and rebroadcasts the checkup message.  This should settle quickly to any major discrepancies.  This consensus may possibly add back dead nodes to the network but will include any live ones known by any node. Thus after local test-live functions are performed on each node to their parent; any dead nodes will be once again weeded out. |

This model maintains a closed loop of node-parent relationships and monitors for failures therein.  A near-real-time model of the current network is maintained at all times and can be made available to an application for service discovery and addressing.  The following describes the process of a node being added to the network:

[Image: Visio Adding to the Network]

**Implementation Challenges**

Liveobjects - "Writing about the limitations of the software environment will also useful. Not in the form of a long and very emotional list of complaints and an opportunity to pour your frustration into it, because those things are best handled by submitting bug reports. But it’s a good place to explain how these limitations affected your project, and what you did to overcome or evade them."

Actual Implementation

Architecture

Features

Configuration (How to Setup)

**Performance Analysis**

Data was gathered on end to end message latency under two sample configurations.  Both tests were performed on a single machine.  The first configuration served as a baseline for the second.  The first used a single transmitter (Source) and single receiver (Viewer).  This was meant to classify performance in a near ideal situation.  Message latency under these conditions was expected to be very low showing the delay caused by the overhead provided by Liveobjects primarily.

Surprisingly this latency was quite high for a localhost test, a simple sockets level piece of software written in c++ reveled latency <1ms while our dataset revealed an average latency of 25ms, a rather significant reduction in performance.  Without the tools to dig deeper our analysis of this issue must stop here, although this overhead must be noted particularly in designing real time applications where greater than 50ms of latency is considered unusable.  With a video distribution system such as ours, buffers can be used to reduce the impact of this lack of performance.  However if this were to be used as the basis for a voip or video conference type piece of software it may be necessary to take more aggressive steps to attain acceptable performance levels.

The second test configuration utilized three separate transmitters (Sources) and three separate receivers (Viewers).  This was expected to result in notably more overhead than the previous example because each viewer would actively filter messages from the sources that were not selected.  As described above a significantly more complex system was initially envisioned which would limit of the impact of these issue.  However this test proved very interesting primarily because it performed at exactly the same level as the previous test.  This would indicate any added latency from inefficiencies in the application were trivial.  Seeing this information was encouraging and in looking forward at continued development would be an excellent indicator as to a strategy to allocate limited development resources.

The interesting transient performance through the first several hundred messages is very interesting.  The settling time is rather long, on the order of minutes, which would seem to rule out hardware or software caching etc.  Without better analysis tools of what the system is doing, it is very difficult to attribute a reason to this performance.  The shape of the transient is notably different in each case, again it is difficult to draw any conclusions from the information available.

Conclusions